# MAR 1 6 2007

## I. AMENDMENTS TO THE SPECIFICATION

Please amend paragraphs [0010], [0031], [0042], and [0043] as indicated below to correct minor typographical errors.

### A. Paragraph [0010]

Please amend sentence two of paragraph [0010] to replace "nodes al-so" with "node also" as shown.

[0010] In another embodiment, a receiving node includes means for receiving and transmitting data frames. The receiving nodes also node also includes means for scheduling nodes to determine the next node to transmit based on a bandwidth allocation and a schedule that specifies a transmit order. The means for scheduling includes means for allowing an out-of-order node to transmit out of the transmit order when the receiver is otherwise idle and means for allowing a new node that is not specified in the bandwidth allocation to transmit when the receiver is otherwise idle.

### B. Paragraph [0031]

Please amend sentence one of paragraph [0031] to replace "208" with "202" as shown.

[0031] In the example shown in FIG. 2, the receiving node 208 202 determines from the received data frame 208 that the sending node 204 has additional data to send (for example, by inspecting a backlog field included in the data frame 208). The receiving node 202 sends a second RTR control frame 210 to the sending node 204. The RTR control frame 210 is addressed to the sending node 204 and invites the sending node 204 to transmit a data frame to the receiving node 202. The sending node 204 receives the RTR control frame 210 and transmits a data frame 212 in response to the RTR control frame 210. Other sending nodes also receive the RTR control frame 210 and determine that the RTR control frame 210 is not addressed to them (for example, by inspecting a destination field included in the control frame 210). As a result, the other sending nodes, at this point, do not transmit and block their own transmissions to avoid interference. The receiving node 202 receives the data frame 212.

#### C. Paragraph [0042]

Please amend sentence one of paragraph [0042] to delete "400" as shown.

[0042] Method 400 is one embodiment of a method 400 of enqueuing data frames that are to be transmitted. Whenever a frame arrives at the MAC layer from a higher layer (checked in block 404), it

is determined if the data frame is a time-critical data frame (checked in block 406). For example, in one embodiment, a predetermined value, for example, is stored by the higher layer in the TTL field of the MAC header of the received data frame to indicate that the received data frame is not a time-critical data frame. If the received data frame is not a time-critical data frame, then the data frame is enqueued for subsequent transmission (block 408).

### D. Paragraph [0043]

Please amend sentence one of paragraph [0043] to insert "of" as shown.

[0044] Method 402 is one embodiment of a method [[of]] dequeing and transmitting a data frame. When a particular node has been invited to transmit data (for example, when a receiving node has sent a RTR control frame to that particular node) (checked in block 420), the sending node selects a data frame for transmission from the queued data frames based on a priority assigned to each of the data frames (block 422). In one such embodiment, the queued data frame with the highest priority is selected for transmission first. Among frames with the same assigned priority, the oldest frame is selected for transmission first. This allows data frames that carry time-critical information access to the medium with less queuing delay, which reduces the total communication latency. In the embodiment shown in FIGS. 4A-4B, a MAC frame header includes an 8-bit priority field that supports up to 256 priority levels. In such an embodiment, the priority is assigned at a higher layer in the OSI protocol, for example, by using a differentiated services QOS protocol such as internet protocol differentiated services (IP DiffServ).